

Statistical Mechanics of Large Dense Random Graphs

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Joint work with Charles Radin

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- A **Large** Random Graph means taking the $n \rightarrow \infty$ limit.
- A Large **Dense** Random Graph means non-trivial edge density.

$$e := \frac{\# \text{ of edges}}{n^2/2} \not\rightarrow 0.$$
$$t := \frac{\# \text{ of triangles}}{n^3/6}.$$

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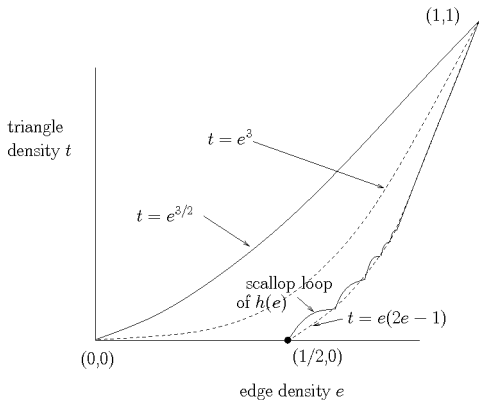
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- What does a typical graph with a given (e, t) look like?

Like microcanonical ensemble with

edges \leftrightarrow particles
triangles \leftrightarrow interaction energy

The shape of (e, t) space



Exponential Random Graphs

Instead of using the microcanonical ensemble, most work has picked parameters β_1, β_2 and set

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Bad idea! (We'll see why soon)

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- 6 Numerical evidence of other phase transitions, some 1st order and some 2nd.

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- A graph G can be viewed as a checkerboard graphon. Break $[0, 1]$ into n equal intervals, each corresponding to a vertex. This breaks $[0, 1]^2$ into n^2 squares, one for each possible edge. Let $g(x, y) = 1$ if the edge exists and 0 if it doesn't.

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- Topology of space of graphons is technical. Buzzword is "Cut metric". Also need to mod out by measure-preserving transformations of $[0, 1]$. (Unnecessary for today's talk.)

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The expected edge and triangle densities are:

$$e(g) = \iint g(x, y) dx dy, \quad t(g) = \iiint g(x, y)g(y, z)g(z, x) dx dy dz.$$

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- Entropy of random process described by graphon is $-n^2 I(g)$
- Entropy density is $-I(g)$.

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The entropy density $s(e, t)$ is well-defined, and equals $-\inf I(g)$, where \inf is over graphons g with $\iint g(x, y) dx dy = e$ and $\iiint g(x, y)g(y, z)g(x, z) dx dy dz = t$.

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Goodbye statistical mechanics. Hello functional analysis.

Two examples

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This is the maximum entropy for fixed e and variable t .

How smooth is the maximum? (Answer: not very)

2) On bottom boundary, minimizing graphon is

$$g(x, y) = \begin{cases} 2e & x < \frac{1}{2} < y \text{ or } y < \frac{1}{2} < x \\ 0 & \text{otherwise.} \end{cases}$$

Entropy is $s(e, 0) = -\frac{1}{2}I_0(2e)$.

The symmetric bipartite phase

Consider the graphon

$$g(x, y) = \begin{cases} e + b & x < \frac{1}{2} < y \text{ or } y < \frac{1}{2} < x \\ e - b & \text{otherwise.} \end{cases}$$

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Rigorous proof that this minimizes $I(g)$ for $e = 1/2$, $t \leq e^3$.
(Later slide)

Theorem (Radin-S)

There exist positive constants c_1, c_2 such that:

$$\text{If } t < e^3, s(e, t) \leq s(e, e^3) - c_1(e^3 - t)^{2/3},$$

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$s(e, t)$ has at most a 1-sided derivative w.r.t. t at $t = e^3$.

$s(e, t)$ is concave **up** for t slightly less than e^3 . Invisible to Legendre transform.

A proof (finally)

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Expand $t = \int \int \int g(x, y)g(y, z)g(z, x)$:

$$t = e^3 + 3e^2 \int \int \delta g(x, y) + 3e \int \int \int \delta g(x, z)\delta g(y, z) + \int \int \int \delta g(x, y)\delta g(y, z)\delta g(z, x).$$

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Last integral is $O(\delta g^3)$, so if $t < e^3$, $e^3 - t < c \|\delta g\|_{L^2}^{3/2}$.

$$s(e, e^3) - s(e, t) = \iint (l_0(e) - l_0(g(x, y))) dx dy = O(\delta g^2).$$

$$s(e, e^3) - s(e, t) = \iint (I_0(e) - I_0(g(x, y))) dx dy = O(\delta g^2).$$

Since

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When $e = 1/2$ and $\delta t < 0$, all estimates are saturated by symmetric bipartite graphon, so we know it's a minimizer.

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Phase II: Asymmetric bipartite. Like symmetric bipartite, but two subintervals have unequal size and unequal probabilities of internal edges. There is a constant $c \neq 1/2$ such that

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Transition from Phase I to Phase II is 2nd-order.

Phase III: Tripartite. Like bipartite, but interval $[c, 1]$ divided into equal sub-intervals. When $x, y > c$, $g(x, y)$ depends on whether x and y are in same or different sub-intervals.

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Transition from Phase II to Phase III is first-order. (Partial derivatives of s diverge).

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Canonical ensemble (e fixed, t variable) misses most of the fun.
Grand canonical ensemble misses almost all of the fun.

Happy Birthday

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Happy Birthday, Yosi!



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